

BASIC COMPRESSION TECHNIQUES

N. C. State University

CSC557 ♦ Multimedia Computing and Networking

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Lectures # 05

Questions / Problems / Announcements?

- Matlab demo of
 - DFT
 - Low-pass windowed-sinc filter design
 - Filtering
- Matlab files are online - try them!
- Compressing the “castle”
- ?

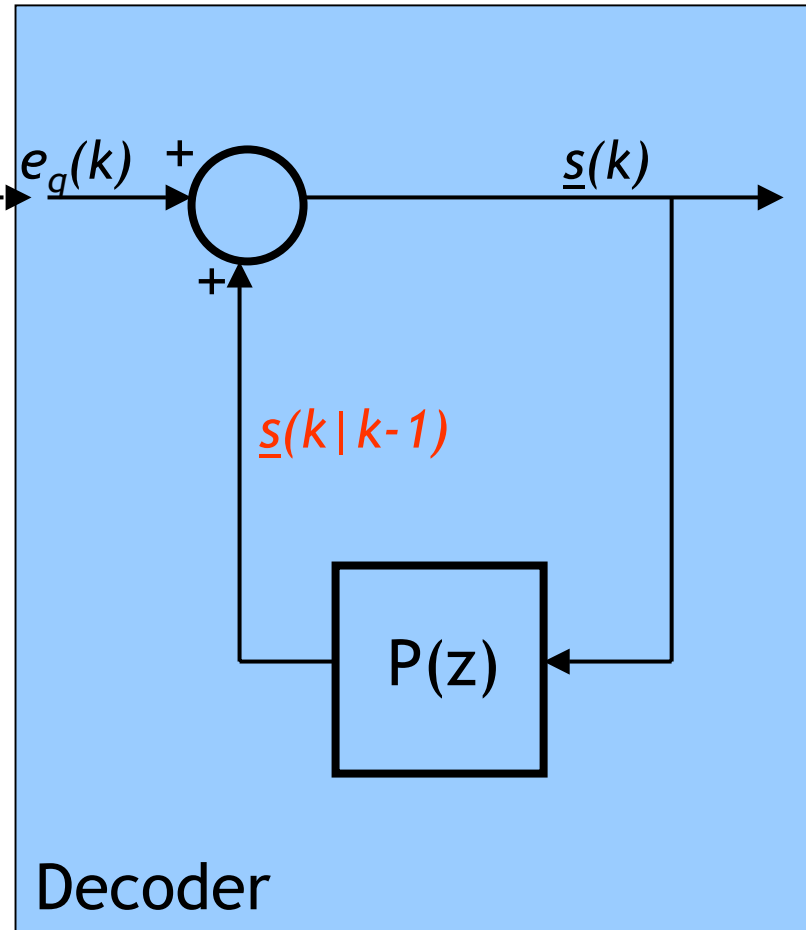
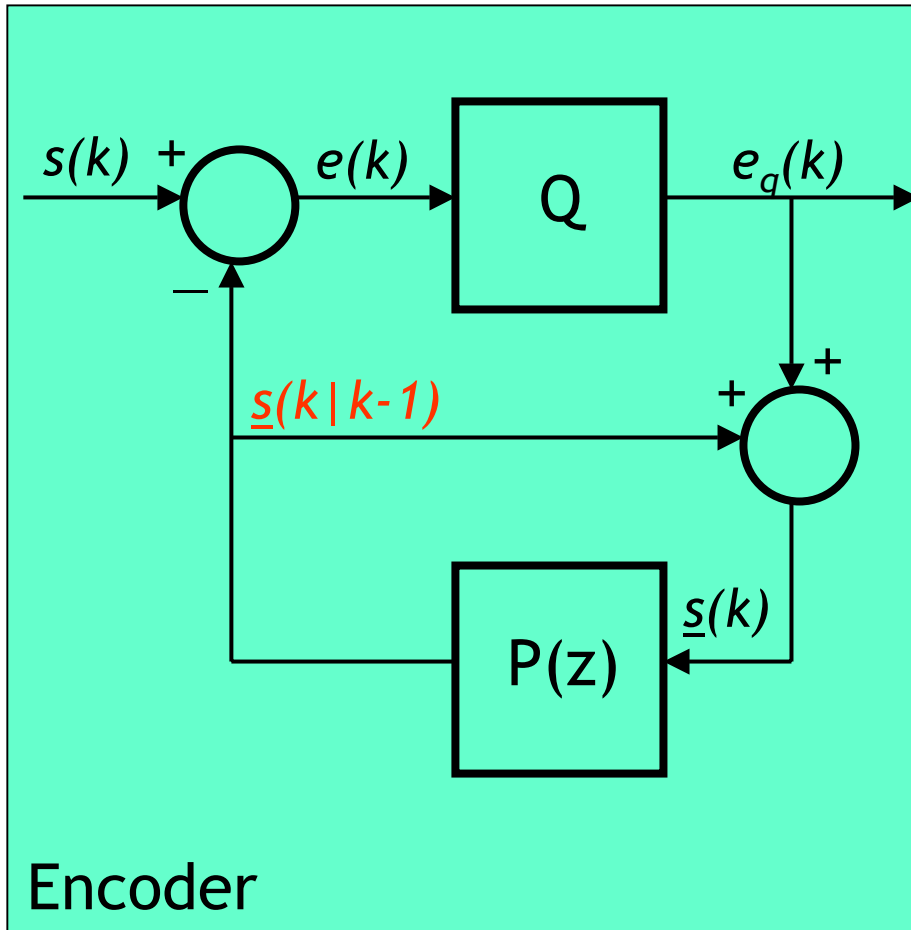
③ Prediction-Based Encoding

- Suppose specific sequences of symbols are predictable, according to some algorithm
- After $i-1$ symbols have been encoded, predict the i^{th} symbol
- Principles involved
 - Redundancy removal \rightarrow subtract a predicted value from the input value
 - *predictor*
 - In other words, successive input samples are correlated
 - Entropy reduction \rightarrow quantize difference between input sample and predicted value
 - *quantizer*

DPCM (cont.)

- Terminology
 - Input signal at time $k = s(k)$
 - Predicted value at time $k = \underline{s}(k|k-1)$
 - Prediction error at time $k = e(k) = s(k) - \underline{s}(k|k-1)$
 - Quantized prediction error at time $k = e_q(k)$
 - Digitized approximation to input at time $k = \underline{s}(k)$
- Quantization
 - Linear PCM, for example
 - Step size may be adjusted, or *adapted*, dynamically
 - (details omitted; see Gibson if interested)

Differential PCM (DPCM)



DPCM Predictor Adaptation

- Prediction accomplished via convolution (filtering)
 - $\underline{s}(k|k-1) = a_1\underline{s}(k-1) + a_2\underline{s}(k-2) + \dots + a_N\underline{s}(k-N)$
- Adapting the prediction: constantly strive to reduce the prediction error
- Example of “backwards adaptation”
 - $a_i(k+1) = \alpha a_i(k) + G(k+1)e_q(k+1)$
 - α is constant “weighting factor” near 1
 - $G(k+1)$ is vector gain computed to minimize the prediction error
 - (details omitted; see Gibson if interested)
- There is no need to transmit the a_i ’s, since the decoder can run exactly the same prediction adaptation algorithm

Example “Media-Specific” Predictors

- Images: the j th Pixel in a row will have the same color as the $(j-1)$ th pixel
- Video: the j th frame will have the same content as the $(j-1)$ th frame

4 LZW Compression

- Given: an alphabet of input “symbols”
 - Pixel colors, ascii characters, ...
- Goal: Construct a "dictionary" of input symbol sequences (ISS's) and output (replacement) symbols
 - desired: output symbols should be shorter than the ISS's they replace
- Initially, the dictionary contains only the individual input symbols
 - “sequences of length 1”
- Scanning input from start to finish
 - add each newly-encountered ISS to the dictionary
 - replace each ISS found in the dictionary with the matching output symbol
 - then *skip over* the matched ISS, and continue

Example (LZW Encoding)

- Assume 12-bit output symbols, 8-bit (ASCII) input symbols
- First 256 entries in dictionary = the alphabet of input symbols
 - 'A'=65, 'B'=66, 'C'=67, ...
- Example of input data to compress: **BABAABAAA**
(=9*8=72 bits)

Encoder Input	Encoder Output	CodeWord Added	String
B	66	256	BA
A	65	257	AB
BA	256	258	BAA
AB	257	259	ABA
A	65	260	AA
AA	260		

LZW Decoding

- Given the encoded (output) data, and the initial dictionary values
 - Do not need the entire dictionary!
- For each output symbol encountered
 - Replace with the input sequence from the dictionary
 - And, create new entry in the dictionary
- If an output symbol is not found in dictionary, it can still be recovered
 - It must be the same as the immediately-previous sequence, with the last symbol repeated

Example (LZW Decoding)

- Encoded string =
— 66 65 256 257 65 260

D e c o d e r I n p u t	D e c o d e r O u t p u t	C o d e W o r d A d d e d	S t r i n g
6 6	B		
6 5	A	2 5 6	B A
2 5 6	B A	2 5 7	A B
2 5 7	A B	2 5 8	B A A
6 5	A	2 5 9	A B A
2 6 0	A A	2 6 0	A A

LZW Dictionary Problems

- Dictionary fills up during construction?
- Options
 - Expand table size dynamically
 - Purge table and start over if compression ratio too low

Lossy Compression Methods

- ⑤ (We've already seen DPCM, ADPCM)
- ⑥ (We'll do transform coding later when we talk about audio and image compression)
- ⑦ Source coding

Source Coders (Synthesis)

- Complex data may be produced by simple mechanism.
Examples:
 - speech from the vocal tract
 - bitmap from a drawing tool
 - musical sound from an instrument
 - image that is a flowchart or block diagram

Source Coders (cont.)

- Compression idea: record the state of the mechanism, instead of the output it produces
- Decompression: *synthesize* the data by modeling the mechanism

Example (Source Coding)

- Compressing a performance on the piano
- Record the keys pressed, when, how long, and how forcefully
- Decompression: synthesize a realistic-sounding piano note
 - with the right pitch
 - at the right time
 - with the right volume
- MIDI recording!

Source Coding Tradeoffs

- Huge amount of compression
- Quality loss because synthesis is not perfect

Sources of Information

- [Crane97] *A Simplified Approach to Image Processing*
 - Chapter 9, section on LZW
- [Gibson98] *Digital Compression for Multimedia*
 - Section 5.3, highlights only