

Link Dimensioning and LSP Optimization for MPLS Networks Supporting DiffServ EF and BE Classes

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Capacity Planning for QoS

- DiffServ + MPLS → QoS in **core** networks
 - DiffServ provides multiple classes of service (e.g., EF, AF, BE)
 - MPLS provides traffic engineering
- Capacity planning is needed
 - **provisioning** adequate resources, **routing** over static LSPs
 - assumes traffic demands predictable in the core network
- **Novel aspect**: capacity planning for multiple classes of service

Challenges

- Tradeoffs between realistic models and solution complexity
- Some aspects of DiffServ not particularly well-defined or mature
 - e.g., AF not included in this work
- Solution of a nonlinear (convex) integer programming problem
 - nonlinear performance constraint, cost function
 - integer routing variables and link capacities

Problem Inputs

- Network topology
- Set of candidate link types (capacities)
 - cost of a link type a non-linear function of capacity
- EF user demands
 - for each O-D (origin-destination) pair, allow m EF demands
 - defined by $\langle \textit{requested bandwidth}, \textit{average arrival rate} \rangle$

Problem Inputs (cont'd)

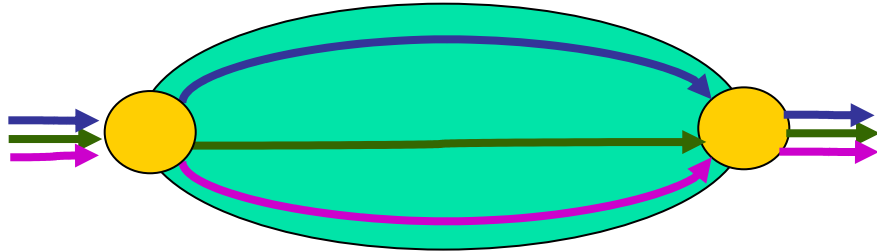
- Best-Effort user demands
 - for each O-D pair, **1** BE demand
 - defined only by *<average arrival rate>*
 - also, define BE "background" traffic demand for each link

Problem Statement

- Objective: minimize the total link cost
- Means:
 - **route** the EF and O-D based BE user demands
 - **choose** the candidate link types, capacities
- Major constraint: BE performance requirement
- Routing constraints:
 - each BE demand **can** be split into multiple paths
 - each EF demand **cannot** be split (in-order delivery guaranteed)

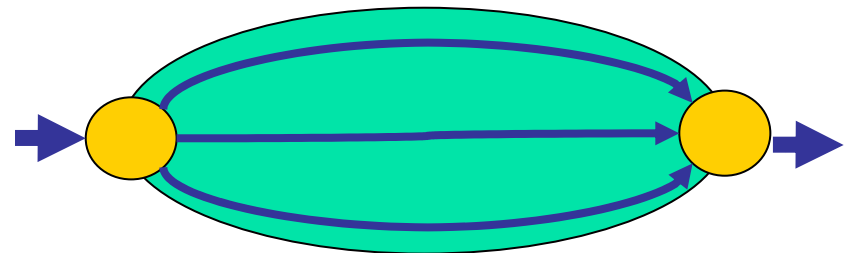
Splitting Traffic

EF user demands



- Multiple EF demands are allowed for the same O-D pair
- Each demand must follow a single LSP

single **BE** user demand



- Only one BE user demand is allowed for the same O-D pair
- Each demand can be split across multiple LSPs

Performance Model

- BE performance requirement is upper bound on average queueing delay **per link**
 - each link = non-preemptive priority queue, M/G/1
 - Poisson arrivals, arbitrary packet length distribution
 - function is convex in the EF arrival rate to the link
- Upper bound = average packet transmission delay * g_i (a user-defined parameter)
- **No** explicit EF performance requirement

Problem Formulation (Alt.)

Goal: $\min \sum_{l \in L} \tilde{C}_l$
 Subject to:

} Total link cost

$$\tilde{\psi}_l \geq f(\beta_{ef}^l)$$

} Performance Constraint

$$x_{kmj}^{ef} = 0/1, \sum_{j \in J_k} x_{kmj}^{ef} = 1$$

} Routing Constraint

$$\sum_{j \in J_k} x_{kj}^{be} = 1$$

$$u_{lt} = 0/1, \sum_{t \in T_l} u_{lt} = 1$$

} Discrete link capacity Constraint

Solution Method

- **Lagrangian relaxation** approach
 - has been successfully applied to many link dimensioning and routing problems
- **Two** sets of discrete variables limit the viable solution method
 - EF routing variables + link sizes
 - relaxing either one of them still results in a difficult non-linear integer programming problem

Solution Details

1. Instead: relax flow aggregation inequalities for both EF and BE traffic

$$\beta_l^{ef} \geq \sum_{k \in K} \sum_{m \in M_k} \alpha_{km}^{ef} \sum_{j \in J_k} x_{kmj}^{ef} \delta_j^l$$

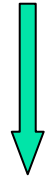
$$\beta_l^{be} \geq \gamma_l + \sum_{k \in K} \alpha_k^{be} \sum_{j \in J_k} x_{kj}^{be} \delta_j^l$$



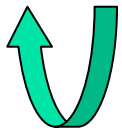
2. Results in three independent (separable) subproblems

- minimize cost per link (solve by gradient projection)
- solve EF routing (shortest path computation)
- solve BE routing (shortest path computation)

Solution Details (cont'd)



3. Update Lagrangean multipliers by Subgradient method



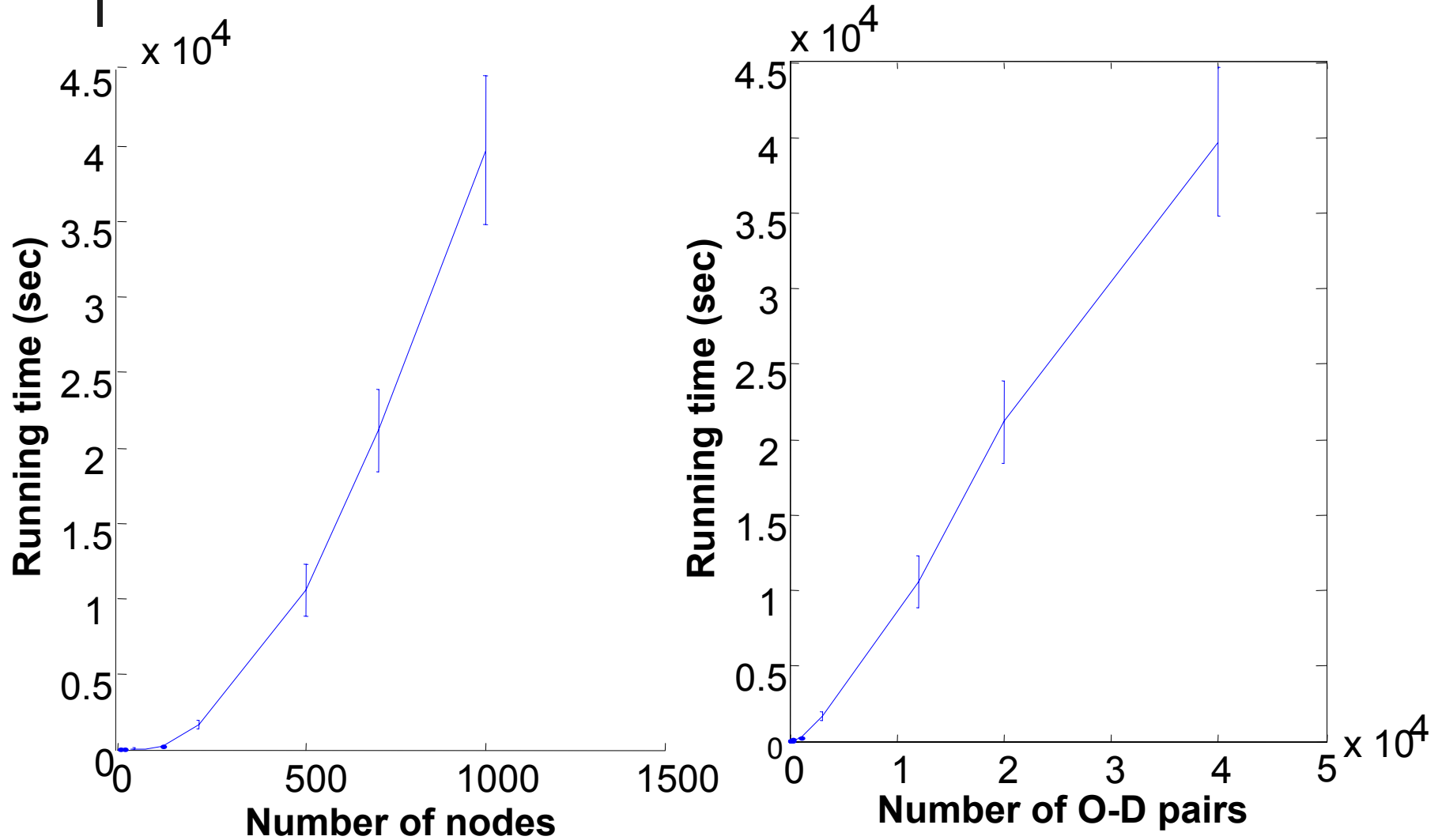
- retain the best solution from all iterations for the primary problem
- maximum of 400 iterations

4. Using best routing solution from dual problem, solve link capacities for primal problem

Computational Results

- Evaluate the algorithm in 8 different sizes of networks
 - up to 1000 nodes and 40000 O-D pairs
 - 30 different topologies for each network size
- Algorithm converges in all 240 test cases
- Measured duality gap $\leq 4\%$
 - upper bound on non-optimality

Computational Results



Contributions

- First capacity planning method for networks using MPLS+DiffServ
- Complex, fairly realistic model
 - BE+EF, priority queueing, O-D and background traffic demands, bifurcated and non-bifurcated routing, discrete set of link types with non-linear costs, ...
- Key solution techniques
 - i. convex performance model
 - ii. novel relaxation

Future Work

1. Incremental (dynamic) solution techniques
2. Include other service classes (AF)
3. Extend to spare capacity planning